

Programming the OpenMP API

Introduction

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2 Programming the OpenMP API Introduction



History

- De-facto standard for Shared-Memory Parallelization.
- 1997: OpenMP 1.0 for FORTRAN
- 1998: OpenMP 1.0 for C and C++
- 1999: OpenMP 1.1 for FORTRAN
- 2000: OpenMP 2.0 for FORTRAN
- 2002: OpenMP 2.0 for C and C++
- 2005: OpenMP 2.5 now includes both programming languages.
- 05/2008: OpenMP 3.0
- 07/2011: OpenMP 3.1
- 07/2013: OpenMP 4.0
- 11/2015: OpenMP 4.5
- 11/2018: OpenMP 5.0
- 11/2020: OpenMP 5.1
- 11/2021: OpenMP 5.2





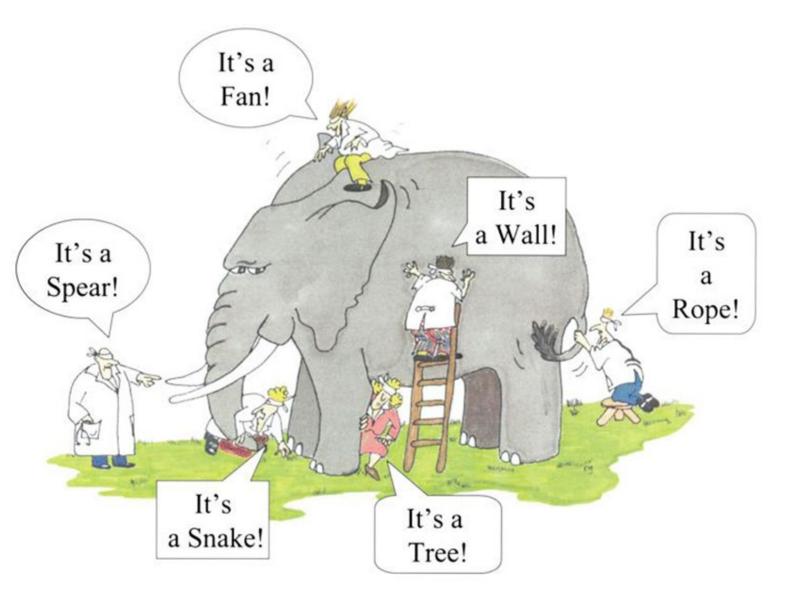
What is OpenMP?

- Parallel Region & Worksharing
- Tasking

•

...

- SIMD / Vectorization
- Accelerator Programming





Get your C/C++ and Fortran Reference Guide! Covers all of OpenMP 5.1/5.2!

OpenMP API 5.1	www.openmp.org	Page 1	
	OpenMP 5.1 API Syntax Referen	ce Guide	
	The OpenMP* API is a portable, scalable model that gives parallel programmers a simple and flexible interface for developing portable parallel applications in C/C++ and Functionality new/changed in OpenMP 5.1 is this oc	Fortran. OpenMP is suitable for a wide range of algorithms running on multicore nodes and chips, NUMA systems, GPUs, and other such devices attached to a CPU.	
May be abbreviated to For	Functionality new/changed in OpenMP 5.1 is this co Functionality new/changed in OpenMP 5.0 in this co		
Directives and Constructs An OpenMP executable directive applies to the successing structured block. A structured block is an OpenMP construct or a block of executable statements with a single entry at			
An OpenMP executable directive applies to the succeeding the top and a single exit at the bottom. OpenMP directive	s except simd and any declarative directive may not appear in	Fortran PURE procedures.	
Variant directives	Informational and utility directives	error [2-5.4] Instructs the compiler or runtime to display a message and	
metadirective [2.3.4] [2.3.4] A directive that can specify multiple directive variants,	requires [2.5.1] [2.4] Specifies the features that an implementation must provide	to perform an error action.	
one of which may be conditionally selected to replace the metadirective based on the enclosing OpenMP context.	in order for the code to compile and to execute correctly.	pragma omp error (clouse ((,) clouse))	
Bpragma omp metadirective [clause] [,] clause]]	#pragma omp requires clouse [[[,] clouse]]	5 Production (1) constants	
+ - or -	5	\$ ISomp error (clouse ([,] clouse)]	
#pragma omp begin metadirective [clause] [,] clause]] ctmt/cl	<pre>§ [Somp requires clause [[[,] clause]]</pre>	clouse:	
Apragma omp end metadirective	clouse: reverse_offload	at(compilation execution) severity(fatal warning)	
<pre>!Somp metadirective [clause] [,] clause]] </pre>	unified address	message(msg-string)	
\$ Somp begin metadirective (clouse((,) clouse))	unified_shared_memory atomic_default_mem_order(seq_cst acq_rel relaxed)	parallel construct	
stmt(s) Somp end metadirective	dynamic_allocators ext_implementation-defined-requirement	parallel [2.6] [2.6]	
clouse: when (context-selector-specification: (directive-variant))	assumes and assume (2.5.2)	Creates a team of OpenMP threads that execute the	
default ((directive-variant/)	Provides invariants to the implementation that may be used	region.	
declare variant (2.3.5) (2.3.5)	for optimization purposes.	5 structured-block	
Declares a specialized variant of a base function and the context in which it is used.	<pre>#pragma omp assumes clouse [[[,] clouse]]</pre>	ISomp parallel (clouse(,,)clouse)) loosely-structured-block	
Pragma omp declare variant(variant-func-id)	 or - #pragma omp begin assumes clouse [[],] clouse]] 	5 onp end parallel	
clause [[[,] clause]]	declaration definition-seq #pragma omp end assumes	tomp parallel [clouse] [,]clouse]]	
(#pragma omp declare variant(variant-func-id) \ clause (([, clause))	- or-	strictly-structured-block (1\$ omp end parallel)	
function definition or declaration	<pre>#pragma omp assume clouse [[[,] clouse]] structured-block</pre>	clowse:	
• or •	Somp assumes clause [[[,] clause]]	default (doto-sharing-attribute) private (list) firstprivate (list) shared (list) copyin (list)	
Apragma omp declare variant clouse declaration-definition-seq	- or -	reduction (/reduction-modifier; / reduction-identifier : /ist) proc_bind (primary master (deprecated) dose spread)	
Apragma omp end declare variant	Somp assume clause [[,] clause]]	allocate (follocator : /list) (C++if (/ parallel :) scolor-expression)	
Somp declare variant ([base-proc name :] & variant-proc name] clause [[[,] clause]]	IS omp end assume	CC++ num, threads (integer-expression)	
clouse:	Somp assume clouse [[[,] clouse]]	For if ([parallel:] scalar-logical-expression) For num_threads (scalar-integer-expression)	
match (context-selector-specification) adjust ares (adjust-op : provment-list)	strictly-structured-block		
adjust_args (adjust-op : argument-list) append_args (append-op[], append-op]]) adjust-ap: nothing, need_device_ptr	clause:	teams construct	
agust-ap: notning, need_device_ptr append-op: interop (interop-type [, interop-type]])	assumption-clause ext_implementation-defined-requirement	teams [2.7] [2.7] Creates a league of initial teams where the initial thread of	
GC++ voriont-func-id: The name of a function variant that	assumption-clause:	each team executes the region.	
is a base language identifier, or for C++, a templote-id. For variant-org-name: The name of a function variant	absent[directive-name [[, directive-name]]) contains[directive-name [[, directive-name]])	#pragma omp teams (clause/ (,)clause)) structured-block	
that is a base language identifier.	no_openmp no_openmp_routines	ISomp teams (clouse) [, iclouse]]	
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Controls whether variant substitution occurs for a given call.	(Ct+ holds(scalar-expression) For holds(scalar-logical-expression)	isomp teams (clouse(), (clouse))	
Ppragma omp dispatch [clouse [[,] clouse]] expression-stmt	nothing (2.5.3)	strictly-structured-block	
Somp dispatch (clause ([,] clause]]	Indicates explicitly that the intent is to have no effect.	/l\$omp end teams/ clouse:	
2 stmt	#pragma omp nothing	private (list) firstprivate (list)	
clouse: depend (/depend-modifier,) dependence type : locator-list)	5	shared (list)	
nowait	15omp nothing	reduction ([default ,] reduction-identifier : list) allocate ([ollocator :] list)	
is_device_ptr(list) C/C++ device (integer-expression)		default (doto-sharing-attribute) num_teams (/ lower-bound :) upper-bound)	
novariants(scolor-expression) (X** nocontext(scolor-expression)		(/C++ thread_limit (integer-expression)	
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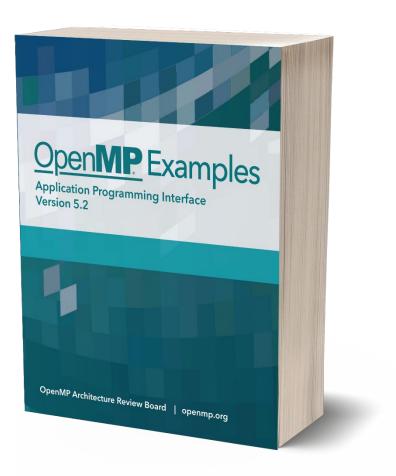
penMP API 5.1	www.openmp.org	Pag
Directives and Constructs (conti	nued)	
hasked construct hasked [24][2:16] ecifies a structured block that is executed by a subset the threads of the current team. [in 5.0, this is the	workshare [2:10.3] [2:8.3] Divides the execution of the enclosed structured block into separate units of work, each executed only once by one thread.	for simd and do simd (2.11.5.2) (2.9.3.2) Specifies that the iterations of associated loops will be executed in parallel by threads in the team and the iterations executed by each thread can also be executed
the interacts of the current ream. In socio di su state construct, un which mader and process maked. Support of the state of the sta	Somp workshare bosity structured block Some of workshare [sowait] - or - Some workshare strictly structured block [Some end workshare [nowait]]	concurrently using SiMD instructions.
Somp masked / filter(scalar-integer-expression) strictly-structured-block [[Somp end masked]	Worksharing-loop construct for and do (2.11.4) [2.9.2]	do directives. declare simd [2115.3] [293.3] Applied to a function or a subroutine to enable the
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	 monotonic: Each thread executes the chunks that is assigned in increasing logical iteration order. A schedule (table) (cluster order cluster implies monotonic. nonemonotonic: Chunks are assigned to threads in any order and the behavior of an application that depends on execution order of the chunks is 	istprivite (it:) collispe (n) dit, checkel (ind, chun; stre)) allocate (officienter:) (int) order (inder-modifer: concurrent) order-modifier: reproducible unconstrained distribute simd (z11.6.3] (2.9.4.3]
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specifies that the associated structured block is executed you have on the threads in the team. Support of the structured block structured Structured Block Structured Structured Structured Structured Structured Structured Structured Structure	Applied to a loop to indicate that the loop can be transformed into a SIMD loop	distribute parallel for and distribute parallel do [2116][20.4] These constructures spectry a loop that can be executed in parallel by multiple threads that are member of multiple topoletic threads and the parallel for [choure[[,choure]] []] []] []] []] []] []] []] []] []] []
[(Somp end single (end_clause[[.]end_clause]]) clause: private [list] firstprivate [list] allocate [(nifoccater:])krt] Cer copyprivate (list] for end_clause: copyprivate [list] nowait	reduction (Induction matifier, I reduction dantifier) lat) collapse (n) order (I (ander modifier) concurrent) dCnif ((and) scalar-legical-expression) For II ((and) scalar-legical-expression) ander-modifier: reproducible unconstrained	clause: Any accepted by the distribute, parallel for,or parallel do directives.
2020 OpenMP ARB		OMP1120-01-ON

OpenMP API Specification & Examples





https://link.openmp.org/book52 https://link.openmp.org/tr11



https://link.openmp.org/examples521

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Recent Books About OpenMP





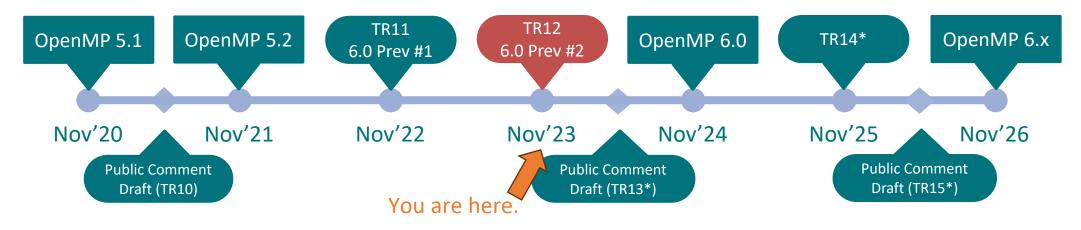
A book that covers all of the OpenMP 4.5 features, 2017 A new book about the OpenMP Common Core, 2019

OpenMP Roadmap



- Roadmap for the releases of the OpenMP API
 - 5-year cadence for major releases, one minor release in between
 - OpenMP 5.2 was an additional release before OpenMP version 6.0
 - (At least) one Technical Report (TR) with feature previews in every year

Version	Year
4.0	2013
5.0	2018
6.0	2024
7.0	2029





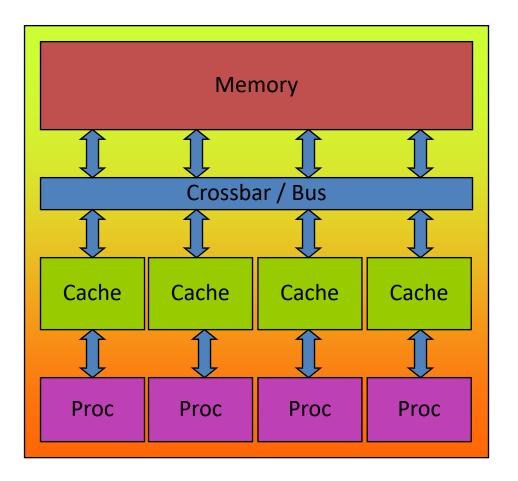
Programming the OpenMP API

Parallel Region

OpenMP's machine model



• OpenMP: Shared-Memory Parallel Programming Model.



All processors/cores access a shared main memory.

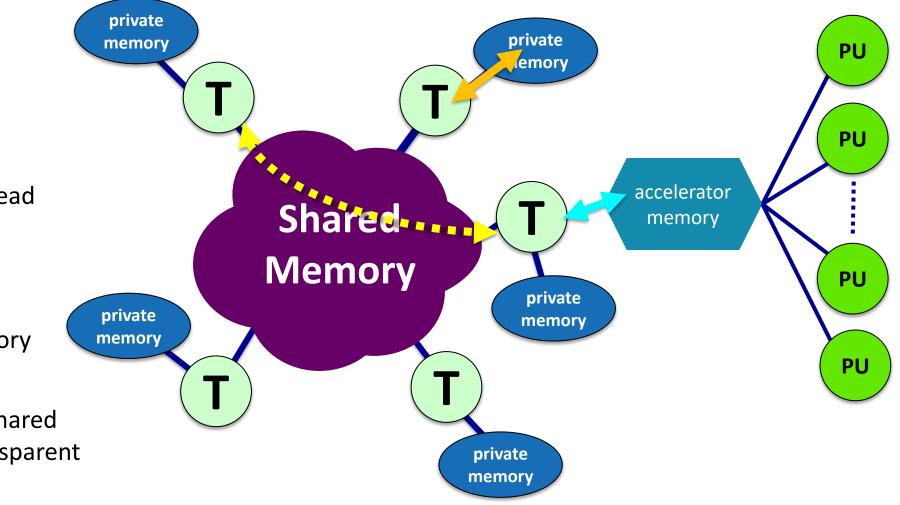
Real architectures are more complex, as we will see later / as we

Parallelization in OpenMP employs multiple threads.



The OpenMP Memory Model

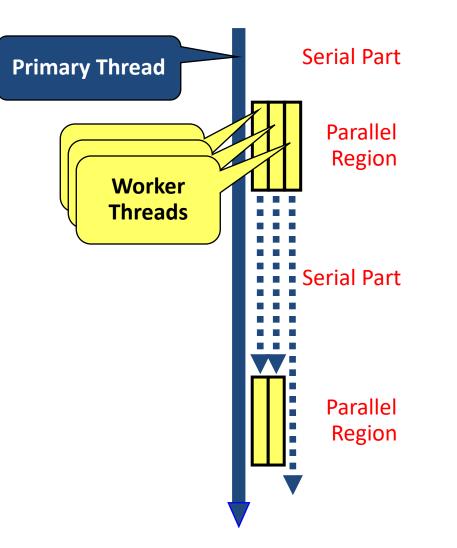
- All threads have access to the same, globally shared memory
- Data in private memory is only accessible by the thread owning this memory
- No other thread sees the change(s) in private memory
- Data transfer is through shared memory and is 100% transparent to the application





The OpenMP Execution Model

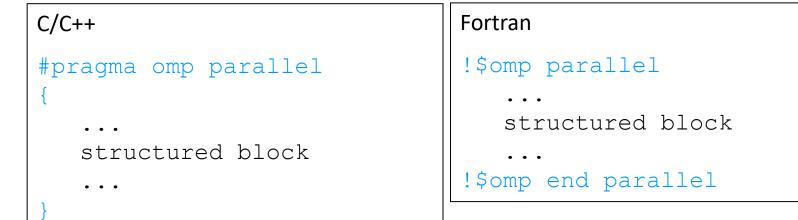
- OpenMP programs start with just one thread: the *Primary Thread*.
- Worker threads are spawned at Parallel Regions, together with the primary thread they form the Team of threads.
- In between Parallel Regions the Worker threads are put to sleep.
 The OpenMP *Runtime* takes care of all thread management work.
- Concept: Fork-Join.
- Allows for an incremental parallelization!





Parallel Region and Structured Blocks

• The parallelism has to be expressed explicitly.



- Structured Block
 - Exactly one entry point at the top
 - Exactly one exit point at the bottom
 - Branching in or out is not allowed
 - Terminating the program is allowed (abort / exit)

- Specification of number of threads:
 - Environment variable: OMP_NUM_THREADS=...
 - Or: Via num_threads clause:
 add num_threads (num) to the parallel construct

OpenMP

Starting OpenMP Programs on Linux

• From within a shell, global setting of the number of threads:

```
export OMP_NUM_THREADS=4
./program
```

• From within a shell, one-time setting of the number of threads:

```
OMP_NUM_THREADS=4 ./program
```



Using OpenMP Compilers

Production Compilers w/ OpenMP Support



GCC

- clang/LLVM
- Intel Classic and Next-gen Compilers
- AOCC, AOMP, ROCmCC
- IBM XL
- ... and many more

See <u>https://www.openmp.org/resources/openmp-compilers-tools/</u> for a list

Compiling OpenMP



- Enable OpenMP via the compiler's command-line switches
 - → GCC: -fopenmp
 - → clang: -fopenmp
 - → Intel: -fopenmp or -qopenmp (classic) or -fiopenmp (next-gen)
 - → AOCC, AOCL, ROCmCC: -fopenmp
 - → HPE/Cray CPE: -homp
 - → IBM XL: -qsmp=omp
- Switches have to be passed to both compiler and linker:

```
$ gcc [...] -fopenmp -o matmul.o -c matmul.c
$ gcc [...] -fopenmp -o matmul matmul.o
$./matmul 1024
Sum of matrix (serial): 134217728.000000, wall time 0.413975, speed-up 1.00
Sum of matrix (parallel): 134217728.000000, wall time 0.092162, speed-up 4.49
```

Demo



Hello OpenMP World



Worksharing



For Worksharing

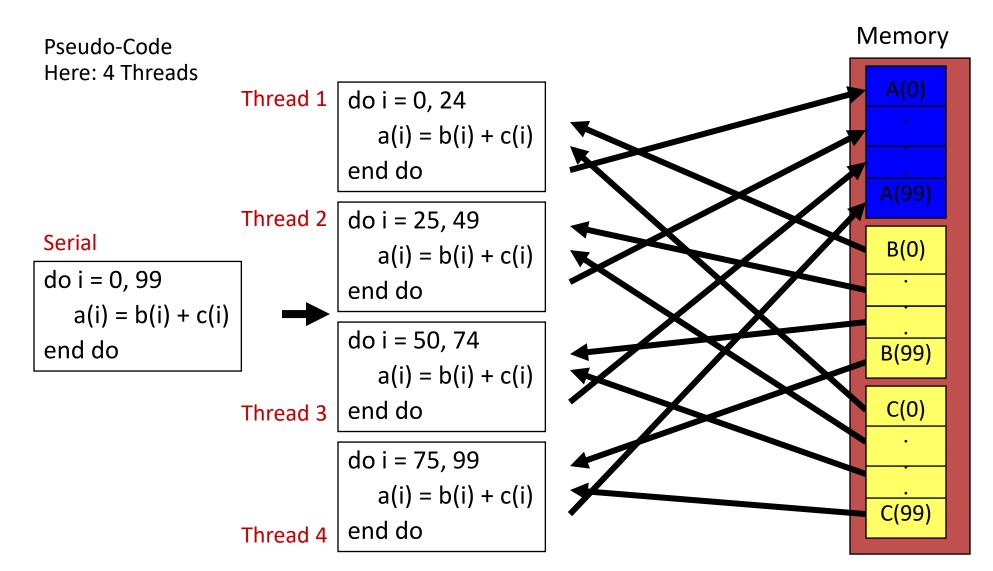
- If only the *parallel* construct is used, each thread executes the Structured Block.
- Program Speedup: *Worksharing*
- OpenMP's most common Worksharing construct: for

C/C++	Fortran
<pre>int i; #pragma omp for for (i = 0; i < 100; i++) {</pre>	<pre>INTEGER :: i !\$omp do DO i = 0, 99 a[i] = b[i] + c[i] END DO</pre>
}	

- Distribution of loop iterations over all threads in a Team.
- Scheduling of the distribution can be influenced.
- Loops often account for most of a program's runtime!

Worksharing illustrated







The Barrier Construct

- **OpenMP** barrier (implicit or explicit)
 - Threads wait until all threads of the current *Team* have reached the barrier

C/C++ #pragma omp barrier

• All worksharing constructs contain an implicit barrier at the end



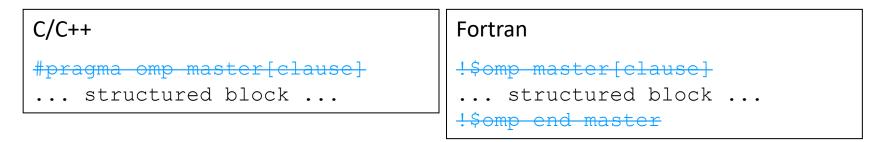
The Single Construct

C/C++	Fortran
<pre>#pragma omp single [clause] structured block</pre>	<pre>!\$omp single [clause] structured block</pre>
	!\$omp end single

- The single construct specifies that the enclosed structured block is executed by only on thread of the team.
 - It is up to the runtime which thread that is.
- Useful for:
 - I/O
 - Memory allocation and deallocation, etc. (in general: setup work)
 - Implementation of the single-creator parallel-executor pattern as we will see later...



The Master Construct is going to be removed with OpenMP 6.0 (2025)



- The master construct specified that the enclosed structured block is executed only by the primary thread of a team.
 - Note: The master construct was no worksharing construct and does not contain an implicit barrier at the end.
- Replacement: see the masked construct later on.

Demo



Vector Addition



Influencing the For Loop Scheduling / 1

- *for*-construct: OpenMP allows to influence how the iterations are scheduled among the threads of the team, via the *schedule* clause:
 - schedule(static [, chunk]): Iteration space divided into blocks of chunk size, blocks are assigned to threads in a round-robin fashion. If chunk is not specified: #threads blocks.
 - schedule(dynamic [, chunk]): Iteration space divided into blocks of chunk (not specified: 1) size,
 blocks are scheduled to threads in the order in which threads finish previous blocks.
 - schedule(guided [, chunk]): Similar to dynamic, but block size starts with implementation-defined value, then is decreased exponentially down to chunk.
- **Default is** schedule (static).

Influencing the For Loop Scheduling / 2



Static Schedule

→ schedule(static [, chunk])

→ Decomposition

depending on chunksize

→ Equal parts of size 'chunksize' distributed in round-robin

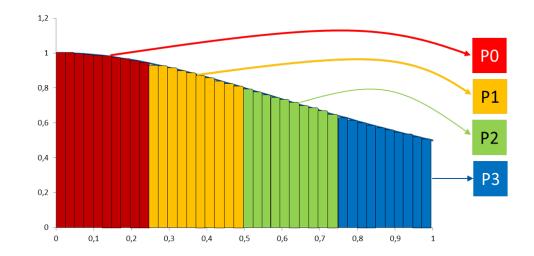
fashion

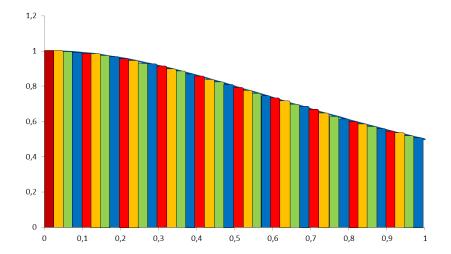
Pros?

→ No/low runtime overhead

Cons?

→ No dynamic workload balancing







Influencing the For Loop Scheduling / 3

- Dynamic schedule
 - schedule(dynamic [, chunk])
 - Iteration space divided into blocks of chunk size
 - Threads request a new block after finishing the previous one
 - Default chunk size is 1
- Pros ?
 - Workload distribution
- Cons?
 - Runtime Overhead
 - Chunk size essential for performance
 - No NUMA optimizations possible



Synchronization Overview

- Can all loops be parallelized with for-constructs? No!
 - Simple test: If the results differ when the code is executed backwards, the loop iterations are not independent.
 BUT: This test alone is not sufficient:

```
C/C++
int i, int s = 0;
#pragma omp parallel for
for (i = 0; i < 100; i++)
{
    s = s + a[i];
}</pre>
```

• *Data Race*: If between two synchronization points at least one thread writes to a memory location from which at least one other thread reads, the result is not deterministic (race condition).



Synchronization: Critical Region

• A Critical Region is executed by all threads, but by only one thread simultaneously (Mutual Exclusion).

C/C++
#pragma omp critical (name)
{
 ... structured block ...
}

• Do you think this solution scales well?

C/C++	
<pre>int i, s = 0; #pragma omp parallel for (i = 0; i < 100; {</pre>	
<pre>#pragma omp critical { s = s + a[i]; }</pre>	}



Scoping

Scoping Rules

OpenMP.

- Managing the Data Environment is the challenge of OpenMP.
- *Scoping* in OpenMP: Dividing variables in *shared* and *private*:
 - *private*-list and *shared*-list on Parallel Region
 - private-list and shared-list on Worksharing constructs
 - General default is *shared* for Parallel Region, *firstprivate* for Tasks.
 - Loop control variables on *for*-constructs are *private*
 - Non-static variables local to Parallel Regions are *private*
 - *private*: A new uninitialized instance is created for the task or each thread executing the construct
 - *firstprivate*: Initialization with the value before encountering the construct
 - *lastprivate*: Value of last loop iteration is written back to the variable in the primary thread
 - Static variables are shared

Tasks are introduced later

Privatization of Global/Static Variables

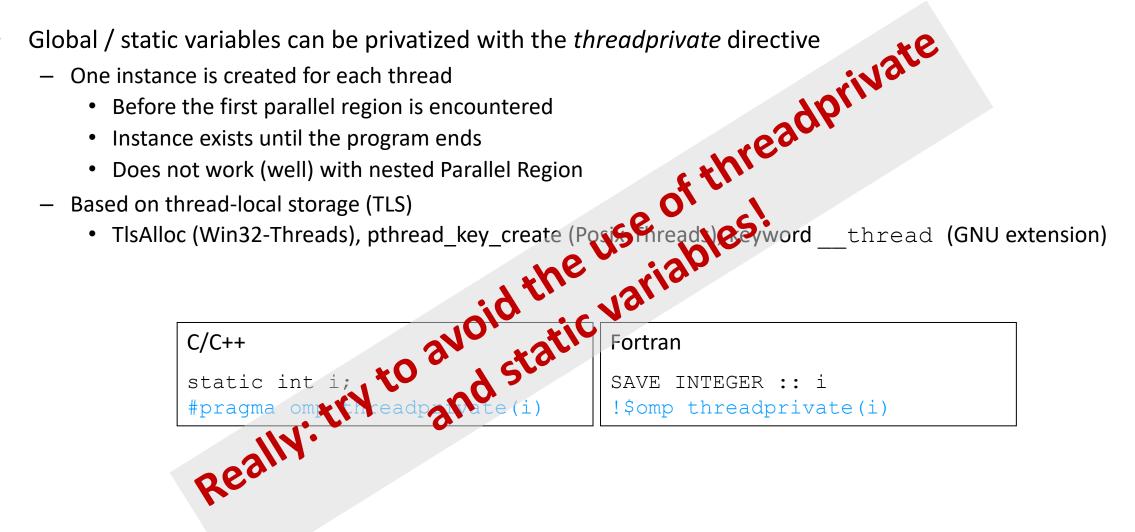


- Global / static variables can be privatized with the *threadprivate* directive
 - One instance is created for each thread
 - Before the first parallel region is encountered
 - Instance exists until the program ends
 - Does not work (well) with nested Parallel Region
 - Based on thread-local storage (TLS)
 - TIsAlloc (Win32-Threads), pthread_key_create (Posix-Threads), keyword ____thread (GNU extension)

C/C++	Fortran
	SAVE INTEGER :: i
<pre>#pragma omp threadprivate(i)</pre>	<pre>!\$omp threadprivate(i)</pre>

Privatization of Global/Static Variables

- ٠





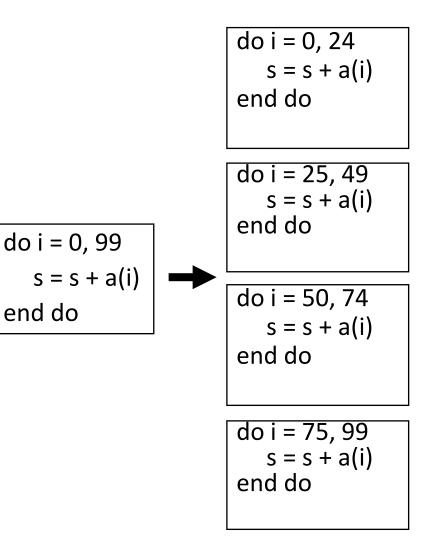
Back to our example

C/C++

```
int i, s = 0;
#pragma omp parallel for
for (i = 0; i < 100; i++)
{
    #pragma omp critical
        { s = s + a[i]; }
}
```

It's your turn: Make It Scale!





#pragma omp parallel

{

}

#pragma omp for
 for (i = 0; i < 99; i++)
 {</pre>

$$s = s + a[i];$$

} // end parallel

(done)



do i = 0, 24

end do

 $S = S + S_1$

end_do

end do

 $S = S + S_{3}$

end do

 $s = s + s_4$

 $s = s + s_{2}$

do i = 25, 49

do i = 50, 74

do i = 75, 99

 $s_1 = s_1 + a(i)$

 $s_2 = s_2 + a(i)$

 $s_3 = s_3 + a(i)$

 $s_4 = s_4 + a(i)$

#pragma omp parallel { double ps = 0.0; // private variable #pragma omp for for (i = 0; i < 99; i++)do i = 0, 99 ps = ps + a[i];s = s + a(i)end do #pragma omp critical s += ps; // end parallel



The Reduction Clause

- In a *reduction*-operation the operator is applied to all variables in the list. The variables have to be *shared*.
 - reduction(operator:list)
 - The result is provided in the associated reduction variable

```
C/C++
int i, s = 0;
#pragma omp parallel for reduction(+:s)
for(i = 0; i < 99; i++)
{
    s = s + a[i];
}</pre>
```

Possible reduction operators with initialization value:

+ (0), * (1), - (0), & (~0), | (0), && (1), || (0), ^ (0), min (largest number), max (least number)

Remark: OpenMP also supports user-defined reductions (not covered here)

Example



Ρ

Example: Pi (1/2)

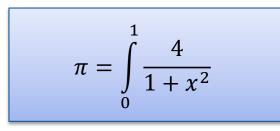
```
double f(double x)
{
    return (4.0 / (1.0 + x*x));
}
```

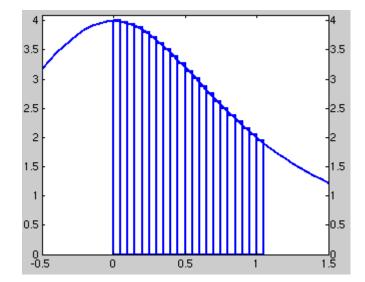
```
double CalcPi (int n)
{
```

```
const double fH = 1.0 / (double) n;
double fSum = 0.0;
double fX;
int i;
```

```
#pragma omp parallel for
for (i = 0; i < n; i++)
{
    fX = fH * ((double)i + 0.5);
    fSum += f(fX);
}
return fH * fSum;
}
```







Example: Pi (2/2)

```
double f(double x)
{
    return (4.0 / (1.0 + x*x));
}
```

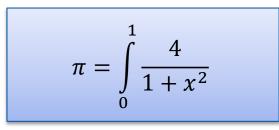
```
double CalcPi (int n)
```

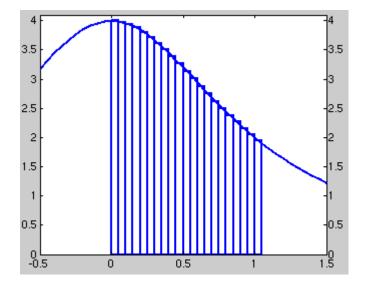
{

```
const double fH = 1.0 / (double) n;
double fSum = 0.0;
double fX;
int i;
```

```
#pragma omp parallel for private(fX,i) reduction(+:fSum)
    for (i = 0; i < n; i++)
    {
        fX = fH * ((double)i + 0.5);
        fSum += f(fX);
    }
    return fH * fSum;
}</pre>
```







Demo



Ρ