

Programming the OpenMP API

Introduction

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History

- De-facto standard for Shared-Memory Parallelization.
- 1997: OpenMP 1.0 for FORTRAN
- 1998: OpenMP 1.0 for C and C++
- 1999: OpenMP 1.1 for FORTRAN
- 2000: OpenMP 2.0 for FORTRAN
- 2002: OpenMP 2.0 for C and C++
- 2005: OpenMP 2.5 now includes both programming languages.
- 05/2008: OpenMP 3.0
- 07/2011: OpenMP 3.1
- 07/2013: OpenMP 4.0
- 11/2015: OpenMP 4.5
- 11/2018: OpenMP 5.0
- 11/2020: OpenMP 5.1
- 11/2021: OpenMP 5.2

What is OpenMP?

- Parallel Region & Worksharing
- Tasking

• …

- SIMD / Vectorization
- Accelerator Programming

Get your C/C++ and Fortran Reference Guide! Covers all of OpenMP 5.1/5.2!

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OpenMP API Specification & Examples

https://link.openmp.org/book52 https://link.openmp.org/tr11

https://link.openmp.org/examples521

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Recent Books About OpenMP

A book that covers all of the OpenMP 4.5 features, 2017

A new book about the OpenMP Common Core, 2019

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OpenMP Roadmap

- Roadmap for the releases of the OpenMP API
	- 5-year cadence for major releases, one minor release in between
	- OpenMP 5.2 was an additional release before OpenMP version 6.0
	- (At least) one Technical Report (TR) with feature previews in every year

Programming the OpenMP API

Parallel Region

OpenMP's machine model

• OpenMP: Shared-Memory Parallel Programming Model.

All processors/cores access a shared main memory.

Real architectures are more complex, as we will see later / as we

Parallelization in OpenMP employs multiple threads.

The OpenMP Memory Model

- All threads have access to the same, globally shared memory
- Data in private memory is only accessible by the thread owning this memory
- No other thread sees the change(s) in private memory
- Data transfer is through shared memory and is 100% transparent to the application

The OpenMP Execution Model

- OpenMP programs start with just one thread: the *Primary Thread*.
- *Worker* threads are spawned at *Parallel Regions*, together with the primary thread they form the *Team* of threads.
- In between Parallel Regions the Worker threads are put to sleep. The OpenMP *Runtime* takes care of all thread management work.
- Concept: *Fork-Join*.
- Allows for an incremental parallelization!

Parallel Region and Structured Blocks

The parallelism has to be expressed explicitly.

- *Structured Block*
	- Exactly one entry point at the top
	- Exactly one exit point at the bottom
	- Branching in or out is not allowed
	- Terminating the program is allowed (abort / exit)
- *Specification of number of threads:*
	- Environment variable: OMP_NUM_THREADS=…
	- Or: Via num_threads clause: add num threads (num) to the parallel construct

Starting OpenMP Programs on Linux

• From within a shell, global setting of the number of threads:

```
export OMP_NUM_THREADS=4
./program
```
• From within a shell, one-time setting of the number of threads:

```
OMP NUM THREADS=4 ./program
```


Using OpenMP Compilers

Production Compilers w/ OpenMP Support

OCC

- clang/LLVM
- Intel Classic and Next-gen Compilers
- AOCC, AOMP, ROCmCC
- **IBM XL**
- … and many more

■ See <https://www.openmp.org/resources/openmp-compilers-tools/> for a list

Compiling OpenMP

- Enable OpenMP via the compiler's command-line switches
	- \rightarrow GCC: -fopenmp
	- \rightarrow clang: -fopenmp
	- \rightarrow Intel: -fopenmp or -qopenmp (classic) or -fiopenmp (next-gen)
	- → AOCC, AOCL, ROCmCC: -fopenmp
	- \rightarrow HPE/Cray CPE: -homp
	- \rightarrow IBM XL: -qsmp=omp
- Switches have to be passed to both compiler and linker:

```
$ gcc [...] -fopenmp -o matmul.o -c matmul.c
 $ gcc [...] -fopenmp -o matmul matmul.o
$./matmul 1024
Sum of matrix (serial): 134217728.000000, wall time 0.413975, speed-up 1.00
Sum of matrix (parallel): 134217728.000000, wall time 0.092162, speed-up 4.49
```
Demo

Hello OpenMP World

Worksharing

For Worksharing

- If only the *parallel* construct is used, each thread executes the Structured Block.
- Program Speedup: *Worksharing*
- OpenMP's most common Worksharing construct: *for*

- Distribution of loop iterations over all threads in a Team.
- Scheduling of the distribution can be influenced.
- Loops often account for most of a program's runtime!

Worksharing illustrated

The Barrier Construct

- OpenMP barrier (implicit or explicit)
	- Threads wait until all threads of the current *Team* have reached the barrier

 $C/C++$ #pragma omp barrier

• All worksharing constructs contain an implicit barrier at the end

The Single Construct

- The single construct specifies that the enclosed structured block is executed by only on thread of the team.
	- It is up to the runtime which thread that is.
- Useful for:
	- I/O
	- Memory allocation and deallocation, etc. (in general: setup work)
	- Implementation of the single-creator parallel-executor pattern as we will see later…

The Master Construct is going to be removed with OpenMP 6.0 (2025)

- The master construct specified that the enclosed structured block is executed only by the primary thread of a team.
	- Note: The master construct was no worksharing construct and does not contain an implicit barrier at the end.
- Replacement: see the masked construct later on.

Demo

Vector Addition

Influencing the For Loop Scheduling / 1

- *for*-construct: OpenMP allows to influence how the iterations are scheduled among the threads of the team, via the *schedule* clause:
	- $-$ schedule (static $\left\lceil$, chunk]): Iteration space divided into blocks of chunk size, blocks are assigned to threads in a round-robin fashion. If chunk is not specified: #threads blocks.
	- schedule(dynamic [, chunk]): Iteration space divided into blocks of chunk (not specified: 1) size, blocks are scheduled to threads in the order in which threads finish previous blocks.
	- $-$ schedule (quided $\left\lceil$, chunk]): Similar to dynamic, but block size starts with implementation-defined value, then is decreased exponentially down to chunk.
- Default is schedule(static).

Influencing the For Loop Scheduling / 2

Static Schedule

- \rightarrow schedule(static [, chunk])
- \rightarrow Decomposition

depending on chunksize

 \rightarrow Equal parts of size 'chunksize' distributed in round-robin

fashion

Pros?

 \rightarrow No/low runtime overhead

Cons?

 \rightarrow No dynamic workload balancing

Influencing the For Loop Scheduling / 3

- Dynamic schedule
	- schedule(dynamic [, chunk])
	- Iteration space divided into blocks of chunk size
	- Threads request a new block after finishing the previous one
	- Default chunk size is 1
- Pros ?
	- Workload distribution
- Cons?
	- Runtime Overhead
	- Chunk size essential for performance
	- No NUMA optimizations possible

Synchronization Overview

- Can all loops be parallelized with for -constructs? No!
	- Simple test: If the results differ when the code is executed backwards, the loop iterations are not independent. BUT: This test alone is not sufficient:

```
C/C++int i, int s = 0;
#pragma omp parallel for
for (i = 0; i < 100; i++){
   s = s + a[i];}
```
• *Data Race*: If between two synchronization points at least one thread writes to a memory location from which at least one other thread reads, the result is not deterministic (race condition).

Synchronization: Critical Region

• A *Critical Region* is executed by all threads, but by only one thread simultaneously (*Mutual Exclusion*).

 $C/C++$ #pragma omp critical (name) { ... structured block ... }

• Do you think this solution scales well?

Scoping

Scoping Rules

- Managing the Data Environment is the challenge of OpenMP.
- *Scoping* in OpenMP: Dividing variables in *shared* and *private*:
	- *private*-list and *shared*-list on Parallel Region
	- *private*-list and *shared*-list on Worksharing constructs
	- General default is *shared* for Parallel Region, *firstprivate* for Tasks.
	- Loop control variables on *for*-constructs are *private*
	- Non-static variables local to Parallel Regions are *private*
	- *private*: A new uninitialized instance is created for the task or each thread executing the construct
		- *firstprivate*: Initialization with the value before encountering the construct
		- *lastprivate*: Value of last loop iteration is written back to the variable in the primary thread
	- Static variables are *shared*

Tasks are

introduced later

Privatization of Global/Static Variables

- Global / static variables can be privatized with the *threadprivate* directive
	- One instance is created for each thread
		- Before the first parallel region is encountered
		- Instance exists until the program ends
		- Does not work (well) with nested Parallel Region
	- Based on thread-local storage (TLS)
		- TlsAlloc (Win32-Threads), pthread_key_create (Posix-Threads), keyword __thread (GNU extension)

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Back to our example

 $C/C++$

int i, $s = 0$; #pragma omp parallel for for $(i = 0; i < 100; i++)$ { #pragma omp critical $\{ s = s + a[i]; \}$ }

It's your turn: Make It Scale!

#pragma omp parallel

{

}

#pragma omp for for $(i = 0; i < 99; i++)$ $\{$

} // end parallel

(done)

 $s_1 = s_1 + a(i)$

 $s_2 = s_2 + a(i)$

 $s_3 = s_3 + a(i)$

 $s_4 = s_4 + a(i)$

#pragma omp parallel { double $ps = 0.0$; // private variable **#pragma omp for** for $(i = 0; i < 99; i++)$ $\{$ $ps = ps + a[i];$ } **#pragma omp critical** { s += $ps;$ } **} // end parallel** do $i = 0, 99$ $s = s + a(i)$ end do do i = $0, 24$ end do $S = S + S_1$ $do i = 25, 49$ end do $S = S + S_2$ do i = $50, 74$ end do $S = S + S_3$ do i = 75, 99 end do $s = s + s_4$

The Reduction Clause

- In a *reduction*-operation the operator is applied to all variables in the list. The variables have to be *shared*.
	- reduction(operator:list)
	- The result is provided in the associated reduction variable

```
C/C++int i, s = 0;
#pragma omp parallel for reduction(+:s)
for(i = 0; i < 99; i^{++})
{
  s = s + a[i];}
```
– Possible reduction operators with initialization value:

+ (0), * (1), - (0), & (~0), | (0), && (1), || (0), ^ (0), min (largest number), max (least number)

– Remark: OpenMP also supports user-defined reductions (not covered here)

Example

PI

Example: Pi (1/2)

```
double f(double x) {
   return (4.0 / (1.0 + x*x)); }
```

```
double CalcPi
(int n)
```
{

```
const double fH = 1.0 / (double) n;
  double fSum = 0.0;
 double fX
;
  int i;
```

```
#pragma omp parallel for
  for (i = 0; i < n; i++){
    fX = fH * ((double)i + 0.5); fSum += f(fX);
   }
 return fH * fSum
;
}
```


Example: Pi (2/2)

```
double f(double x)
{
   return (4.0 / (1.0 + x*x));
}
```

```
double CalcPi (int n)
```
{

```
const double fH = 1.0 / (double) n;
 double fSum = 0.0;
 double fX;
 int i;
```

```
#pragma omp parallel for private(fX,i) reduction(+:fSum)
  for (i = 0; i < n; i++) {
    fX = fH * ((double)i + 0.5); fSum += f(fX);
 }
   return fH * fSum;
}
```


Demo

